

#### **396EM Airline Operations and Scheduling / 6075MAA Airline Scheduling and Operations**

#### Lecture 1c Airline Network Flows and Integer Programming Models

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# **Integer Linear Programming**



 A large part of the problems that airlines face can be translated into network and integer programming models.

- In short, an integer programming problem is a mathematical optimization or feasibility program in which some or all of the variables are restricted to be integers
- We provide a review of some of the optimization models



## **Network Flow models**



1. Shortest Path (Router) problem

#### 2. Minimum cost flow problem

#### 3. Maximum flow problem

#### 4. Multi-commodity Problem



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#### Network





#### **Basic elements of a network**





# Network Terminology

- Nodes and Arcs: in a network, the points (circles) are called nodes and the lines are referred to as arcs, links, or arrows.
- Flow: the amount of goods, flights, passengers, and etc. that move from one node to another.
- Directed Arc: if the flow through an arc is allowed only in one direction, then the arc is said to be a directed arc.
- Undirected Arc: When the flow on an arc can be move in either direction, it is called an undirected arc.













# Supply Node and Demand Node

- Arc Capacity: the maximum amount of flow that can be sent through an arc, e.g. restrictions on the flight no. between 2 cities.
- Supply nodes: Nodes with the amount of flow coming to them greater than the amount of flow leaving them, or nodes with positive net flow.
- Demand Nodes: Nodes with negative net flow or outflow greater than inflow.
- Transshipment: Nodes with the same amount of flow arriving and leaving or nodes with zero net flow.
- Path: a path is a sequence of distinct arcs that connect 2 nodes in this fashion since sometimes 2 nodes are not connected by an arc, but could be connected by a sequence of arcs. E.g. Airlines utilize hubs to provide connections between city pairs in their network.





# Cycle and Connected Network

Source: Starting node in the path.

- Destination: Last node in the path.
- Cycle: A sequence of directed arcs that begins and ends at the same node. E.g. Aircraft start from an airport which is a maintenance base and, after flying to several destinations, end up at the same airport from which they departed.
- Connected Network: A network in which every two nodes are linked by at least one path.







#### **Network Flow models**







#### 1. Shortest Path (Route) problem



- Attempts to identify a path, from source to destination, within the network, that results in minimum transport time/cost.
- The objective is to identify the path with the minimum cost between two nodes.
- The problem consists of a connected network with known costs for each arc in the network
- Typical application for air cargo handlers, freight forwarders and etc. to identify the shortest route to transport cargos through indirect flights



#### Notation:



#### Network with flight times between city pairs

- The nodes represent the cities, and the arcs are the flights.
- The numbers on the arcs represent the flight time in minutes between the city pairs.
- We want to determine the best route that results in the shortest flying time from node





# Example: Network with flight times between city pairs





## Solution



- Decision Variable: ?
- Objective function: ?
- Constraints: ?





#### **Decision Variable**



We assume the following binary (0-1) decision variable:

$$x_{i,j} = \begin{cases} 1 \text{ if arc } (i,j) \text{ is part of the solution} \\ 0 \text{ otherwise} \end{cases}$$



# **Objective Function**



p.14

Minimise:

 $70x_{1,2} + 63x_{1,3} + 56x_{1,4} + \dots + 72x_{7,10} + 87x_{5,10} + 97x_{6,10} + 69x_{9,10}$ 



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### Constraints



Source node (Origin node): the flow must originate from node 1 and it must be a part of the solution, therefore:

 $x_{1,2} + x_{1,3} + x_{1,4} = 1$ 







p.16

# Constraints (Con't)

Source node (Origin node):

 $x_{1,2} + x_{1,3} + x_{1,4} = 1$ 

Same as Destination node: the flow must end up at the destination node (node 10).

 $x_{5,10} + x_{6,10} + x_{7,10} + x_{9,10} = 1$ 



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#### **Constraints for Transshipment Nodes**

- Transshipment nodes: every other node (except origin and destination) is a transshipment node. The net flow in these nodes should be zero since good are not left there, they must be sent to the final destinations.
- The constraints for those transshipment nodes

$$x_{1,2} + x_{4,2} + x_{3,2} - x_{2,3} - x_{2,4} - x_{2,5} - x_{2,6} - x_{2,7} = 0$$

$$x_{1,3} + x_{2,3} + x_{4,3} - x_{3,2} - x_{3,4} - x_{3,5} - x_{3,6} = 0$$

$$x_{1,4} + x_{2,4} + x_{3,4} - x_{4,2} - x_{4,3} - x_{4,5} - x_{4,6} - x_{4,9} = 0$$

$$x_{2,5} + \dots = 0$$

$$x_{2,6} + \dots = 0$$

$$x_{2,7} + \dots = 0$$

$$x_{5,8} + \dots = 0$$

$$x_{4,9} + \dots = 0$$

p.17

## **Final solution**



- The minimum cost is 198
- The shortest path is 1-4-6-10





#### Summary



Sets M = Set of nodes

- Index i, j, k = Index for nodes
- Parameters

C<sub>i, j</sub> = Cost of flow along the arc joining node *i* to node *j m* = Destination node







#### Decision Variable

 $x_{i,j} = \begin{cases} 1 \text{ if arc } (i,j) \text{ is part of the path} \\ 0 \text{ otherwise} \end{cases}$ 

Objective Function

$$Minimize \sum_{i \in M} \sum_{j \in M} \sum_{i,j \in M} x_{i,j}$$

Subject to

$$\sum_{j \in M} x_{1,j} = 1 \quad j \neq 1$$

$$\sum_{j \in M} x_{i,j} - \sum_{k \in M} x_{k,i} = 0 \quad For \ all \ (\forall) \ i, \ i \neq 1 \ and \ i \neq m$$

$$\sum_{i \in M} x_{i,m} = 1$$

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## **Network Flow models**







## 2. Minimum cost flow problem



- Seek to satisfy the requirements of nodes at minimum cost.
- It is a generalized form of transportation, transshipment, and shortest path problems.
- The problem assume the cost per unit of flow and capacities associated with each arc are known.



# Air cargo problem description

- An airline is tasked with transporting goods from nodes 1 and 2 to nodes 5, 6, 7 in the diagram.
- The airline does not have direct flights from the source nodes to destination nodes, but they have the network which can connect those nodes through their hubs in nodes 3 & 4.
- The numbers next to the nodes represent the demands/supply in tons.
- Objective: determine the best way to transport the goods from sources to destinations to minimise the total cost.
- Constraint: The aircraft flying to and from node 4 can carry a maximum of 50 tons of cargo







#### Network presentation for minimum cost flow



Note: If we carry 75 tons of cargo from node 1 to node 3, the cost is 5\*75=375





#### Solution: To formulate this cargo problem

- Decision variable: x<sub>i,i</sub> = amount of flow from node i to node j
- The objective function is:

Minimize:  $5x_{1,3} + 8x_{1,4} + 7x_{2,3} + ...$ 

The constraints:  $x_{1,3} + x_{1,4} \le 75$   $x_{2,3} + x_{2,4} \le 75$   $x_{1,3} + x_{2,3} - x_{3,5} - x_{3,6} - x_{3,7} = 0$   $x_{1,4} + x_{2,4} - x_{4,5} - x_{4,6} - x_{4,7} = 0$   $x_{3,5} + x_{4,5} = 50$   $x_{3,6} + x_{4,6} = 60$  $x_{3,7} + x_{4,7} = 40$   $75 \underbrace{1}_{8}^{5} \underbrace{3}_{5}^{1} \underbrace{5}_{7}^{5} \underbrace{50}_{7}^{7} \underbrace{3}_{7}^{7} \underbrace{3}_{4}^{7} \underbrace{3}_{4}^{7} \underbrace{6}_{6}^{6} \underbrace{60}_{7} \underbrace{7}_{7} \underbrace{2}_{4} \underbrace{4}_{4}^{7} \underbrace{7}_{4} \underbrace{7}_$ 



Additional constraints about capacity

All the flights to and from node 4 (Airport Y) can carry a maximum of 50 tons (might because of the shorter length of its runaway)

$$x_{1,4} \le 50$$
  
 $x_{2,4} \le 50$   
 $x_{4,5} \le 50$   
 $x_{4,6} \le 50$   
 $x_{4,7} \le 50$ 





## **Final Solution: Air cargo**



Total Minimum cost of \$1,250.





## **Network Flow models**







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## 3. Maximum Flow Problem



#### A Special Case of Minimum Cost Flow Problem

- E.g. An airline must determine the number of daily connecting flights that can be arranged between Daytona Beach (DAB), Florida, and Lafayette (LAF), Indiana.
- Connecting flights must stop in Atlanta (ATL), Georgia, and then make one more stop in either Chicago(ORD), Illinois, or Detroit (DTW), Michigan.



#### Example:



#### Maximum number of flights per city-pair

The airline wants to determine how to maximize the number of connecting flights daily from Daytona Beach, FL, to Lafayette, IN, respecting the current restrictions.

City-Pairs	Maximum number of daily flights
DAB - ATL	3
ATL - ORD	2
ATL - DTW	3
ORD - LAF	1
DTW - LAF	2



#### Solution



Work out the network presentation







#### Decision Variable

- Xij = Number of flights (integer from node i to node j)
- f = Number of daily flights from DAB to LAF
- Maximize daily flights between DAB and LAF
   => Maximize f
- Source Node

DAB is the source node, f is the total flow leaving DAB  $x_{1,2} = f$ 





Transshipment nodes (example node 2 ATL)

• 
$$x_{1,2} - x_{2,3} - x_{2,4} = 0$$

- Similarly, transshipment constraints for other nodes 3 and 4.
- Destination node
  - same number of daily flights f departing from DAB should now arrive at destination node LAF.

• 
$$x_{4,5} + x_{3,5} = f$$





- Arc capacity
- The last set of constraints address the capacity of arcs as follows:
  - $x_{1,2} \le 3$  $x_{2,3} \le 2$  $x_{2,4} \le 3$  $x_{3,5} \le 1$  $x_{4,5} \le 2$



# **Final Solution**



- Solving this problem generates a maximum flow of three daily flights between DAB and LAF as follows:
  - 1 flight assigned to the DAB-ATL-ORD-LAF route, and;
  - 2 flights assigned to the DAB-ATL-DTW-LAF route.



### Summary



#### Sets

M = Set of nodes

Index

i,j,k =Index for nodes

#### Parameters

$L_{ii}$	= Lower bound on flow through arc $(i,j)$
Ů,,	= Upper bound on flow through arc $(i,j)$
m	= Destination node

#### Decision Variables:



= Amount of flow from node *i* to node *j*= Amount of flow from source node to destination node



# Summary (Con't)



#### Maximize f

#### Subject to

$$\sum_{j \in M} x_{1, j} = f \qquad \Leftrightarrow \text{ Origin Node}$$

$$\sum_{i \in M} x_{i,j} - \sum_{k \in M} x_{j,k} = 0 \quad \Leftrightarrow \text{Transshipment nodes}$$

$$\sum_{i \in M} x_{i,m} = f$$

$$L_{i,j} \le x_{i,j} \le U_{i,j}$$

 $\Leftrightarrow$  Destination node



## **Network Flow models**







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# 4. Multi-Commodity Problem



- All the network models explained so far assume that a single commodity or type of entity is sent through a network.
  - A network may transport different types of commodities.
- The multi-commodity problem seeks to minimize the total cost when different types of goods are sent through the same network.



#### Example



We modify the example that was presented for the Minimum Cost Flow problem discussed earlier to address the multi-commodity model formulation.





## Solution



#### Decision Variable

- To formulate this problem we assume the following decision variable.
- In this decision variable the indices i and j represent the nodes (i,j = 1,.,7) and k represents the type of commodity (k = 1,2).
- x i,j,k = Amount of flow from node i to node j for commodity k

#### Objective Function

• Minimize  $5x_{1,3,1} + 5_{x1,3,2} + 8x_{1,4,1} + 8x_{1,4,2} + \dots$ 





We need to write one constraint for each node. For example, for node 1 we have:

$$\begin{aligned} x_{1,3,1} &+ x_{1,4,1} \le 40 \\ x_{1,3,2} &+ x_{1,4,2} \le 35 \end{aligned}$$

Similar constraint for the other six nodes





All the flights to and from node 4 can carry a maximum of 50 tons

$$\begin{aligned} x_{1,4,1} + x_{1,4,2} &\leq 50 \\ x_{2,4,1} + x_{2,4,2} &\leq 50 \\ x_{4,5,1} + x_{4,5,2} &\leq 50 \\ x_{4,6,1} + x_{4,6,2} &\leq 50 \\ x_{4,7,1} + x_{4,7,2} &\leq 50 \end{aligned}$$



## **Final Solution**



Solving this problem using software generates a total minimum cost of \$1,250.





## **Key References**



- M. Bazargan (2010) Airline Operations and Scheduling.
   2nd edition, Ashgate
  - Chapter 2 Network Flows and Integer Programming Models

